

A BxtendDSL Solution to the TTC2023 Asymmetric and Directed Bidirectional Transformation for Container Orchestrations Case

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Abstract

Container orchestration is a critical component in the realm of DevOps practices, enabling the efficient management of containers within complex application architectures. However, a significant challenge arises in reconciling the disparity between high-level graphical representations of container orchestration models and the concrete configuration files essential for container orchestration tools. To address this issue, this paper proposes a novel bidirectional and asymmetric transformation approach, facilitating the translation from container orchestrations to MiniYAML through the utilization of BxtendDSL, a hybrid bidirectional and incremental model-to-model transformation language capable of supporting both declarative and imperative specification of model transformations. The paper presents the proposed solution, outlining the transformation rules, and assesses the effectiveness of the approach using benchmark criteria.

Keywords

container orchestration, bidirectional transformations, model merging, graphical models, YAML, BxtendDSL

1. Introduction

The transformation case described in this paper is of significant relevance to DevOps engineers and addresses real-world scenarios. Leite et al. [1] define DevOps as a collaborative and multidisciplinary effort within organizations to automate the continuous delivery of new software versions while ensuring correctness and reliability. With the increasing interest in DevOps, various domain-specific modeling notations have been created, covering aspects like microservice architectures [2], DevOps processes [3], and multi-cloud applications [4].

DevOps heavily relies on tools that facilitate automated deployment, often accomplished by reading declarative descriptions written in structured formats like YAML¹. These structured formats adhere to loosely defined schemas that may evolve across different versions, exemplified by the Docker Compose file format².

The proposed case draws inspiration from Piedade et al. [5] and revolves around container orchestration using Docker Compose. It entails transforming an Ecore [6] model representing the abstract syntax of a high-level graphical DSL for container orchestration into another Ecore model representing the abstract syntax of a Docker Compose YAML file. The transformation must be bidirectional, allowing changes in both models to be prop-

agated back and forth seamlessly. Additionally, changes in the high-level graphical DSL should be conveyed to the YAML file while preserving elements containing information that cannot be represented in the high-level DSL.

This transformation case results in a directed but asymmetrical transformation process. In this paper, we present our solution to this proposed transformation using BxtendDSL [7, 8, 9], our hybrid language designed for bidirectional and incremental model transformations.

The paper is structured as follows: In Section 2, we provide an overview about BxtendDSL. Section 3 describes both the declarative and imperative parts of our solution to the transformation case, followed by a detailed evaluation according to different criteria in Section 4. Section 5 concludes the paper.

2. BxtendDSL

BxtendDSL [7, 8, 9] is a state-based framework for defining and executing bidirectional incremental model transformations on demand that is based on EMF [6] and the programming language Xtend³. It builds upon Bxtend [10], a framework that follows a pragmatic approach to programming bidirectional transformations, with a special emphasis on problems encountered in the practical application of existing bidirectional transformation languages and tools.

When working with the stand-alone Bxtend framework, the transformation developer needs to specify both transformation directions separately, resulting in Bxtend

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¹<https://yaml.org/>

²<https://docs.docker.com/compose/compose-file/>

³<https://eclipse.dev/Xtext/xtend/>

transformation rules with a significant portion of repetitive code.

To this end, BXtendDSL adds a declarative layer on top of the BXtend framework, which significantly reduces the effort required by the transformation developer. Figure 1 depicts the layered approach of our tool: First, the external DSL (BXtendDSL Declarative) is used to specify correspondences declaratively. Second, the internal DSL (BXtendDSL Imperative) is employed to add algorithmic details of the transformation that can not be expressed on the declarative layer adequately.

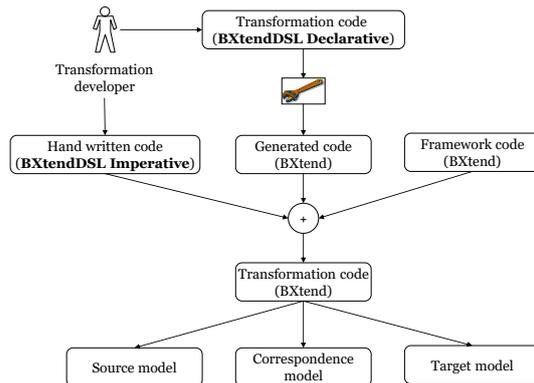


Figure 1: Layered approach used in BXtendDSL

The handwritten code and the generated code are combined with framework code to provide for an executable transformation. The transformation developer is relieved from writing repetitive routine parts of the transformation manually using a code generator. The generated code ensures round-trip properties for simple parts of the transformation. Since the declarative DSL usually is not expressive enough to solve the transformation problem at hand completely, the generated code must be combined with handwritten imperative code. Certain language constructs of the declarative DSL define the interface between the declarative and the imperative parts of the transformation. From these constructs, *hook methods* are generated, the bodies of which must be manually implemented. Hook methods are used, e.g. for implementing filters or actions to be executed in response to the deletion or creation of objects, etc.

Incremental change propagation relies on a persistently stored *correspondence model*, which allows for $m : n$ correspondences between source and target model elements. A powerful *internal DSL* may be used at the imperative level, to retrieve correspondence model elements associated with a given element from the source and target models, respectively. Please note that the transformation developer does not have to deal with managing correspondences at the declarative level, rather all the al-

gorithmic details of managing the correspondence model are handled by our framework automatically.

3. Solution

In this section, we explain the details of our BXtendDSL solution for the Asymmetric and Directed Bidirectional Transformation for Container Orchestrations. We will discuss the different layers in separate subsections. The sources for our solution are publicly available on Github: [GitHub](#)

3.1. Declarative Layer

BXtendDSL code at the declarative layer is used to define transformation rules between elements of source and target models respectively. Listing 1 depicts the code for the transformation at the declarative layer.

```

1 sourcemodel "http://york.ac.uk/ttc/containers/1.0.0"
2 targetmodel "http://york.ac.uk/ttc/miniyaml/1.0.0"
3
4 rule Volume2MapEntry
5 src Volume v;
6 trg MapEntry me | filter;
7
8 v.name <--> me.key;
9
10 rule Image2MapEntry
11 src Image img;
12 trg MapEntry me | filter, creation;
13
14 img.image <--> me.value;
15
16 rule VolumeMount2Scalar
17 src VolumeMount vm;
18 trg Scalar sc | filter;
19
20 vm.path vm.volume --> sc.value;
21
22 rule Container2MapEntry
23 src Container c;
24 trg MapEntry me | filter;
25
26 c.name <--> me.key;
27 c.image c.replicas c.dependsOn {c.volumeMounts:
28   VolumeMount2Scalar} --> me.value {me.value:
29   VolumeMount2Scalar};
30 c.image c.replicas c.dependsOn <-- me.value;
31
32 rule Composition2Map
33 src Composition c;
34 trg Map m | filter, creation;
35
36 {c.nodes: Image2MapEntry, Container2MapEntry, Volume2MapEntry}
37 --> m.entries {m.entries: Image2MapEntry,
38   Container2MapEntry, Volume2MapEntry};
39 c.nodes <-- m.entries {m.entries: Image2MapEntry[img],
40   Container2MapEntry[c], Volume2MapEntry[v]};
  
```

Listing 1: BXtendDSL code at the declarative layer

The code at the declarative layer comprises the transformation rules for all required model elements. Each rule consists of *src* and *trg* patterns. The *trg* patterns contain modifiers, such as *filter* and *creation*, which result in the generation of hook methods. The implementation of the hook methods is described in Section 3.2. After declaring the *src* and *trg* patterns in each rule, the respective

mappings are declared. Mappings may be bidirectional, depicted by the double arrow <-->, or specified for a certain transformation direction, for example, forward (-->) or backward (<--). A very simple bidirectional mapping is depicted in Line 8 of Listing 1: the attribute name of a Volume is assigned to the attribute key of a MapEntry and vice versa.

Note that BxtendDSL was intentionally left incomplete when designed to keep the language as simple and as small as possible. To this end, we did not incorporate an expression language into BxtendDSL. Instead, we decided to apply the generation gap pattern [11] and generate hook methods that are called from the generated code at the respective steps during the transformation. Bodies for hook methods are supplied at the imperative layer using the Xtend programming language.

Consequently, a mapping that has one element on each side of the arrow may be transformed directly into executable code. Hook methods are generated for mappings with more than one element on one side of the arrow symbol. The mapping depicted in line 20 of Listing 1 is used to map the attributes path and volume to the attribute value of the target pattern. Because the declarative language does not comprise mechanisms to describe how the two attributes are mapped to a single attribute on the target side, a hook method is generated (cf., Listing 2).

The transformation specification on the declarative layer also comprises mappings of (containment) references between the elements of the source and target models. Lines 27 and 28 depict the respective mapping in the forward and backward directions. In the forward direction (Line 27), the cross reference image, attribute replicas, cross reference dependsOn, and containment reference volumeMounts of the class Container are mapped to the respective MapEntry in the YAML model. Please note that Bxtend rules are executed in their textual order specified in the BxtendDSL declarative file. That is, rules Volume2MapEntry, Image2MapEntry, and VolumeMount2Scalar are executed before rule Container2MapEntry. Thus, we can be sure that all the elements required for this mapping actually exist and can be retrieved. The syntax of the mapping specified in line 27 contains curly brackets. This indicates that for this feature, the correspondence/trace model is accessed to obtain the respective model elements from the source and target models. The specification of this mapping also results in the generation of a hook method, which is used to describe all the algorithmic details to realize this mapping on the imperative layer.

Rule Composition2Map (c.f., Line 30-35 in Listing 1) maps the root elements of both models. These elements (transitively) contain all other model elements via references nodes and entries. To realize the transformation, this rule is executed after all other rules are executed to ensure that the respective model elements actually exist

when they are assigned to the containment references of the source and target root elements. In the forward direction, Images, Containers, and Volumes are assigned to the respective entries of the target Map. Again, please note that curly brackets are used to access the correspondence model and retrieve the already existing target elements for Images, Containers, and Volumes. A similar mapping is required for the backward transformation (see Line 35 in Listing 1).

3.2. Imperative Layer

On the imperative layer, the bodies for hook methods must be supplied. This holds for the specification of modifiers (e.g., filter or creation), as well as for mappings where further information is required, which cannot be supplied using the declarative language.

Listing 2: Hook method for mapping attributes path and volume to Scalar.value

```
1  override protected valueFrom(String path, Volume
    volume) {
2      return new Type4value(volume.name + ":" + path)
3  }
```

Listing 2 shows the imperative code that is required to realize the mapping vm.path vm.volume --> sc.value, as depicted in Line 20 of Listing 1. The value attribute of the scalar is a concatenation of the name of the Volume and the path, separated by ":".

Please note that this rule does not specify the backward direction; rather, it is addressed in the imperative code for mapping c.nodes <-- m.entries ... from the rule Composition2Map in Listing 1.

Listing 3 depicts the code required on the imperative layer to realize the rule Image2MapEntry, as specified in Lines 10-14 in Listing 1.

Listing 3: Imperative code for rule Image2MapEntry

```
1  override protected filterMe(MapEntry me) {
2      me.key == "image"
3  }
4
5  override protected onMeCreation(MapEntry me) {
6      me.key = "image"
7  }
8
9  override protected valueFrom(String image) {
10     new Type4value(MiniyamlFactory.eINSTANCE.
        createScalar() =>
11         [value = image])
12 }
13
14 override protected imageFrom(Value value) {
15     return new Type4image(((value as Scalar).value))
16 }
```

The implementation of the modifiers filter and creation is shown in Lines 1-7 of Listing 3. The modifiers result in the generation of methods filterMe and onMeCreation, and the transformation developer only needs to supply a body to realize the desired behavior. In this case, an Image from the container model is transformed into a MapEntry of the YAML model. The corresponding MapEntry has a key attribute with the value "image". The filter is applied when transforming in backward direction and it is used to filter all MapEntry elements from the target model and only retrieve the ones whose key attribute contains the value "image".

Methods valueFrom and imageFrom are hook methods that are generated from the mapping depicted in Line 14 of Listing 1. Note that in this case, the mapping only contains a single element on each side of the arrow, but the respective attributes are of different types. Thus, hook methods are required to specify how these types are mapped onto each other. In the forward direction (method valueFrom), the "image" String is transformed into a Scalar, where the attribute value is assigned to the value attribute of the scalar. The value of this attribute is returned in the backward direction.

Listing 4: Imperative code for mapping specified in Line 27 of Listing 1

```

1  override protected valueFrom(Image image, int
      replicas,
2  List<Container> dependsOn, List<Scalar> volSc,
3  Value oldValue) {
4      var entry = yamlFactory.createMap()
5      if (replicas > 1) {
6          val me = yamlFactory.createMapEntry() =>
7              [key = "replicas"
8               value = yamlFactory.createScalar() =>
9                 [value = "" + replicas]
10             ]
11         entry.entries += me
12     }
13     if (image != null)
14         entry.entries += (elementsToCorr.get(image).
15             getTarget()
16             .get(0) as SingleElem).element as MapEntry
17         if (!dependsOn.isEmpty) {
18             val me = createMapEntry("depends_on")
19
20             val list = me.value as miniyaml.List
21             for (Container c : dependsOn)
22                 list.values += yamlFactory.createScalar() =>
23                     [s | s.value = c.name]
24             entry.entries += me
25         }
26     if (!volSc.empty) {
27         val me = createMapEntry("volumes")
28         val list = me.value as miniyaml.List
29         for (Scalar s : volSc) {
30             list.values += s
31         }
32         entry.entries += me
33     }
34     ...

```

```

33     new Type4value(entry)

```

Listing 4 depicts the imperative code that is required to realize the mapping of image, replicas, dependsOn and volumeMounts attributes and references of a Container to respective MapEntries in forward direction. The imperative code contains several conditional blocks that need to be processed if certain conditions hold, for example, if the value of the integer attribute replicas is > 1. In this case, a new MapEntry with appropriate key-value pairs is created and added to the parent Map. If additional MapEntries are required for dependencies and volumes, they are also created using this hook method. Respective entries are then also added to the map, which is then returned at the end of the hook method.

4. Evaluation

The aim of the proposed transformation case is to answer research questions concerning conciseness (i.e., how much specification effort is required to solve this case with current bx tools), preservation of information that cannot be mapped on the other model, and scalability (i.e., how well the proposed solution scales with increasing model sizes).

To this end, the transformation is classified according to the evaluation criteria discussed in the following subsections.

4.1. Correctness

The correctness of the transformation may be verified using two different and supplied comparators. The *MiniYAMLComparator* ignores the order of elements in the respective models, whereas the *MiniYAMLExactComparator* also considers the ordering of elements.

MiniYAMLComparator The BXtendDSL solution passes all supplied tests (forward, backward, incremental forward) for the benchmarx testsuite using the *MiniYAMLComparator*; for example, we achieved 100% correctness in this case.

MiniYAMLExactComparator For the *MiniYAMLExactComparator*, the BXtendDSL solution passes eight out of nine batch forward tests, seven of eight batch backward tests, and four of five incremental forward tests, which leads to a correctness rate of 86.4%. The following test cases fail, due to incorrect ordering of elements in a multi-valued reference: completeModel in both BatchForward and BatchBackward, as well as updateReplicas in IncrementalForward

4.2. Conciseness

To measure the conciseness of the transformation specification, nodes in the respective AST/ASG of the languages used should be counted. To this end, solution developers are required to provide specific implementations of the AST-counter. In our case, an AST-counter for the BxtendDSL language and an AST-counter for the Xtend programming language are required. However, an additional problem remains. First, Bxtend does not provide static libraries; rather, framework code is generated specifically for each transformation. That is, the project contains a significant portion of the generated Xtend and Java code, which must be excluded when the AST nodes are computed. Furthermore, large parts of the Xtend code that are used to implement hook methods on the imperative layer are generated as well. Consequently, the transformation developer is required to specify only the bodies of the respective languages. However, an AST-counter works on a valid source code, that is, an Xtend class that has no compile errors and counts all nodes present in the class. This would lead to incorrect results because large parts of the generated code would be considered. From our understanding, however, conciseness should only take the parts of the code into account that the transformation developer must supply to make the transformation work. To this end, we decided to provide conciseness information using the LOC metrics defined in [12] and [13]. We further split up the numbers into code required on the declarative and the imperative layer respectively. Table 1 presents the results.

	BxtendDSL Declarative	BxtendDSL Imperative
Lines of code	32	202
Number of words	100	788
Number of characters	862	5967

Table 1
Size of the transformation definitions of both solutions

The results clearly indicate that a significant portion of the transformation was specified on the imperative layer. This is due to the asymmetric nature of the transformation case, which cannot be handled adequately in the declarative layer of BxtendDSL. However, the resulting transformation specification is still concise if we compare it to similar transformation cases such as AST2Dag [14].

4.3. Least Change

Beyond correctness, the transformation should preserve additional information in the YAML file that cannot be expressed in the Containers model. When executing the transformation in the two modes (*MiniYAMLComparator*, which checks the general correctness, and *MiniYAMLEx-*

actComparator, which also considers the order of elements in the YAML model), we observe that the transformation can preserve the additional information specified in the YAML file, but not in the exact order of elements. Test cases specified in class *IncrementalForward* were considered for this test. While the BxtendDSL solution achieves 100% accuracy for Mode 2, it passes four out of five tests for Mode 1, resulting in an accuracy of 80% when the exact order of the elements matters.

4.4. Scalability

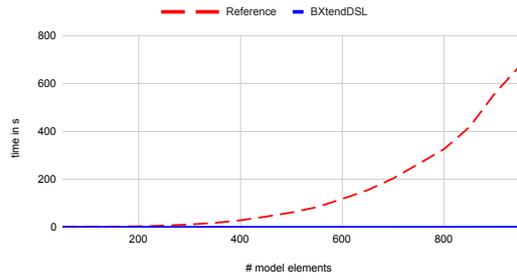
To measure how well the transformation scales to models with increasing numbers of containers, volumes, and images, different scalability tests have been executed in the forward and backward directions, as well as both in batch and incremental situations. We attempted to compare the results of BxtendDSL to the provided reference implementation, but unfortunately, the backward measurements failed with errors on our computers. To this end, only the plots for batch forward and incremental forward transformations contain data for the reference implementation as well. All scalability tests were performed on the same computer in isolation to avoid side effects. A desktop PC with an AMD Ryzen 7 3700x CPU was used, running at a standard clock of 3.60 GHz, with 32 GB of DDR4 RAM and Microsoft Windows 11 64-bit as the operating system. We used Java 13.0.2, Eclipse 4.27.0, and EMF version 2.33.0, to compile and execute Java code for the scalability test suite. Each test was repeated five times, and the median measured time was computed.

For each test, we used the provided class for scalability measurements, which created models of increasing sizes up to 1000 elements. BxtendDSL proves to scale very well with increasing model sizes, as depicted by the plots for batch forward (c.f., Figure 2), incremental forward (c.f., Figure 3), batch backward (c.f., Figure 4), and incremental backward (c.f., Figure 5). Please note that two plots are given per figure: one with linear scaling of the x and y axes and the other with logarithmic scaling. The linear plot is meant to provide a realistic impression for the actual complexity curve of the BxtendDSL solution compared to the reference implementation. The logarithmic plots help zoom into finer details for smaller models (practically invisible in the linear plot), and zoom out for larger models so even large differences in runtime can still be presented qualitatively.

5. Conclusion

The BxtendDSL solution provided for the asymmetric and directed bidirectional transformation for the container orchestration case has proven to be concise, sufficiently correct, and scalable.

Batch Forward: Reference Implementation and BxtendDSL



Batch Forward: Reference Implementation and BxtendDSL

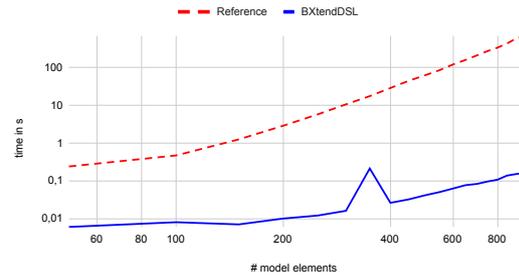
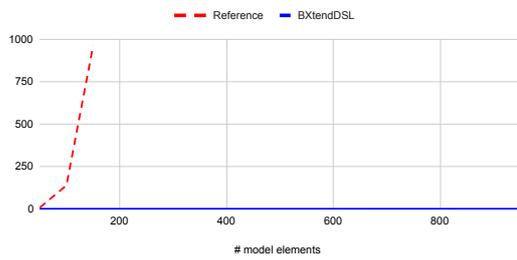


Figure 2: Forward batch transformation: Linear/linear scale (left) and log/log scale (right)

Incremental Forward: Reference Implementation and BxtendDSL



Incremental Forward: Reference Implementation and BxtendDSL

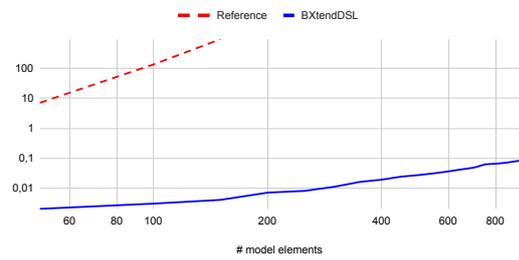


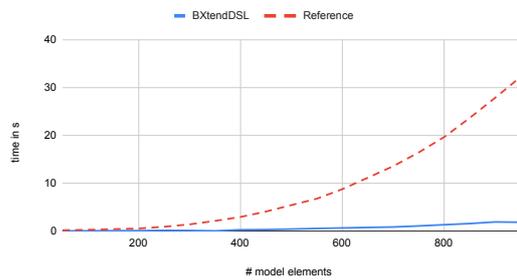
Figure 3: Forward incremental transformation: Linear/linear scale (left) and log/log (right)

Because BxtendDSL allows us to specify details of the transformation on both declarative and imperative levels, the transformation developer may choose (almost) freely which programming paradigm is best suited for the transformation problem at hand. A combination of both results in high expressive power while simultaneously maintaining low specification effort at the same time.

The transformation case revealed small bugs in the

code generation engine, which was used to generate executable code from declarative specifications. Thus, minor tweaks of the generated code are required. These issues have already been addressed, and will be incorporated into BxtendDSL in the upcoming release.

Batch Backward: Reference Implementation and BxtendDSL



Batch Backward: Reference Implementation and BxtendDSL

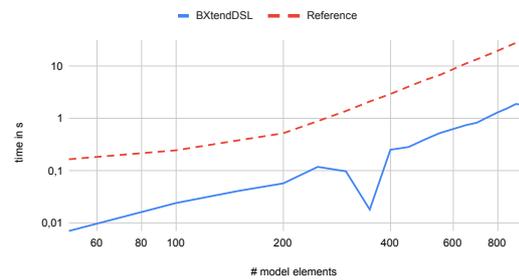
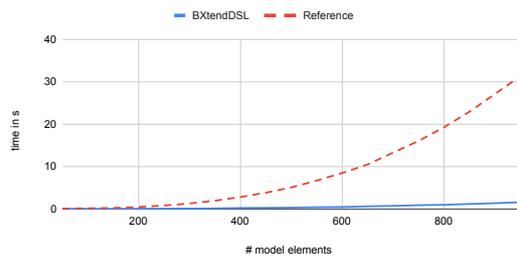


Figure 4: Backward batch transformation: Linear/linear scale (left) and log/log (right)

Incremental Backward: Reference Implementation and BxtendDSL



Incremental Backward: Reference Implementation and BxtendDSL

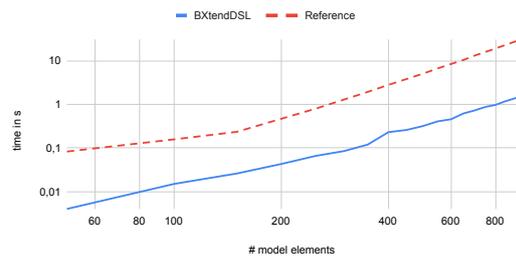


Figure 5: Backward incremental transformation: Linear/linear scale (left) and log/log (right)

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